

ELLIOTT

903

Volume 1: FUNCTIONAL SPECIFICATION
Part 2: THE BASIC 903 COMPUTER UNIT
Section 2: WORD FORMAT AND INSTRUCTION CODE

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Chapter 1: INTRODUCTION

1.1 General.

The 903 is a conventional stored program computer operating in the parallel binary mode. In the computer each word of information consists of an 18-bit binary pattern which may represent either an instruction or an operand of an instruction (e. g. a constant).

The accumulator (A-register) usually holds the result of executing an instruction, and this result is then available as one of the operands for the next instruction. For some purposes - chiefly during the arithmetic functions of multiplication and division - the less significant end of the accumulator is extended by the addition of the 17 most significant bits of the Q register, to hold a 35-bit operand.

Instructions are stored and obeyed sequentially and without gaps. A number of registers apart from A may be used while instructions are being obeyed. Some registers are accessible to program and they are dealt with fully in Section 1. 2. 1; those not accessible to program are dealt briefly in Section 1. 2. 1. and explained fully in Part 1 of Volume 4.

Chapter 2: WORD FORMAT

2.1 General.

The binary bits of a word are referred to as bits 18, 17, 16..... 3, 2, 1, bit 1 being the least significant.

2.2 Representation of Numbers.

When a word represents a number, bit 18 is the sign bit and indicates whether the number following is positive or negative; the binary point is placed immediately after bit 18. Numbers are represented as fractions and when bit 18 is zero the following 17 bits represent a positive number; when it is 1 a negative number is represented. Positive numbers are represented by the appropriate digits of bits 17-1 being set to 1, and negative numbers are represented inversely in "two's complement" form.

The value that may be attributed to each digit of a word is as follows:

Bit no.	18, 17,	16.....	3,	2,	1
Value	-1,	2^{-1} ,	2^{-2}	2^{-15} ,	2^{-16} , 2^{-17}

Thus the largest positive number that may be represented is:

$$0.11\ 111\ 111\ 111\ 111\ 111\ 111\ (1 - 2^{-17})$$

The smallest positive number is:

$$0.00\ 000\ 000\ 000\ 000\ 001\ (2^{-17})$$

The negative number with the largest possible modulus is:

$$1.00\ 000\ 000\ 000\ 000\ 000\ (-1)$$

The negative number with the smallest possible modulus is:

$$1.11\ 111\ 111\ 111\ 111\ 111\ (-2^{-17})$$

2.3 Representation of Instructions.

Instructions are of the single address type and when a word represents an instruction, its bits are grouped as follows:-

Group symbol	B	F	N
Bits per group	1	4	13

Chapter 3: INSTRUCTION CODE

3.1 Instruction Table.

NOTES: (1) The following symbols are used in this table:

p = number of places shifted

z = number of words transferred

dr = device response time

For significance of all other symbols used in this table see Introduction to the 903 Manual.

- (2) All times specified are for unmodified instructions and are subject to a $\pm 10\%$ tolerance. For a modified instruction $7.2 \mu\text{s}$ must be added to the times specified.
- (3) The obeying of any instruction in these tables means that $s:=s+1$ unless otherwise stated.
- (4) If an instruction is to be modified, q is altered in an undefined manner before the instruction is obeyed.

FUNCTION	DESCRIPTION	EFFECT OF INSTRUCTION	REGISTERS NOT AFFECTED	INSTRUCTION TIMES (μ S)	NOTES
0	Set B - Register	b: = m q: = m	A	29.5	an INTERRUPT cannot take place after this instruction.
1	Add	a: = a+m	Q,B	23.0	
2	Negate and Add	a: = -a+m q: = m	B	26.0	
3	Store Auxiliary Register	m18: = 0 m17-1: =q18-2	A,Q,B	24.5 24.5	
4	Load	a: = m	Q,B	23.0	
5	Store	m: = a	A,Q,B	24.5	
6	Collate	a: = a & m	Q,B	23.0	This function generates the logical product of a and m
7	Jump if a zero	Provided a = 0, s13-1: = n s16-14 unchanged s18, 17 undefined q undefined If a \neq 0, s:=s+1	A,B	If a=0, 27.5 If a>0, 21.0 If a<0, 19.5	
8	Jump unconditionally	Whether a is negative, zero or positive (not tested), the effect is as for Function 7, except that s:=m	A,Q,B	23.0	
9	Jump if a -ve	Provided a<0, the effect is as for Function 7. If a \geq 0, s:=s+1	A,B	If a<0, 25.0 If a \geq 0, 19.5	
10	Count in store	m: = m+2 ⁻¹⁷	A,Q,B	23.5	
11	Store SCR	q18-17 undefined q16-14:=s16-14 q13-1:=0 m13-1:=s13-1 m18-14:=0	A,B	30.0	q18, q17 are undefined because s18, s17 are undefined
12	Multiply	(a, q18-2):=axm q1 undefined	B	78.5	

FUNCTION	DESCRIPTION	EFFECT OF INSTRUCTION	REGISTERS NOT AFFECTED	INSTRUCTION TIMES (μ s)	NOTES
13	Divide	$a := \frac{a, q18-2}{n} + 2^{-17}$ $q := \frac{a, q18-2}{n} - 2^{-17} + 2^{-17}$ $a1 := 1$ $q1 := 0$	B	79.0	<p>It is not in general possible to say when the result in A is greater or less than the true quotient. If, however, the quotient can be expressed exactly in 17 or fewer bits (counting from the sign digit) then the following alternatives apply:-</p> <p>a) Divisor positive. The correct quotient is in Q; A contains the correct quotient plus 2^{-17}</p> <p>b) Divisor negative. The correct quotient is $A + 2^{-17}$; the contents of Q are the correct quotient less 2^{-16}.</p>
14	Shift				
	Left Shift	Provided $0 \leq n \leq 47$ $(a, q) := (a, q) \times 2^n$	B		The effect of trying to shift more than 48 places is not defined.
	Right Shift	Provided $8144 \leq n \leq 8191$ $(a, q) := (a, q) \times 2^{n-8192}$	B		The sign bit is regenerated.
	Block transfer Block input	Provided $2048 \leq n \leq 4095$ Transfer x words from the peripheral device specified by bits n11-1 into store locations y to y+x-1 $y = a$ $x = q12-1$ (i.e. $x \leq 4095$) If x=0 the instruction has no effect.	B	23.5+(7+dr)z	Before the instruction is issued A must contain the address of the first location to which data is to be input. q12-1 must specify the number of locations to which data is to be input. On completion of the instruction, A contains the last word input.
Block output	Provided $4096 \leq n \leq 6143$ Transfer x words to the peripheral device specified by bits n11-1 from store locations y to y+x-1 $y = a$ $x = q12-1$ (i.e. $x \leq 4095$) If x=0 the instruction has no effect.	B		Before the instruction is issued A must contain the address of the first location from which data is to be output. q12-1 must specify the number of locations from which data is to be output. One completion of the instruction, A contains the last word output.	

FUNCTION	DESCRIPTION	EFFECT OF INSTRUCTION	REGISTERS NOT AFFECTED	INSTRUCTION TIMES (μs)	NOTES
15	Input Input information from peripheral device	Provided $0 \leq n \leq 2047$ a _n = one 18-bit word from the device specified by bits n11-1	Q,B	20.5+dr	
	Input from Tape Reader	Provided $n=2048$, a shifted left 7 places a8-1 = 8-bit character input from tape reader	Q,B		If both a1=0 before left shift takes place, and bit 8 of character read by tape reader = 0, a8=0; if a1 was a one, bit 8 of the character read by tape reader = 1
	Input from Teleprinter	Provided $n = 2052$, a shifted left 7 places a8-1 = 8-bit character input from teleprinter	Q,B		m16-14 are ignored
	Output Output information to peripheral device	Provided $4096 \leq n \leq 6143$ one 18-bit word output from A to the device specified by bits n11-1	Q,B	19.0+dr	
	Output information to Tape Punch	Provided $n=6144$ a8-1 output to paper tape punch	Q,B		
	Output information to Teleprinter	Provided $n=6148$, a 8-1 output to teleprinter	Q,B		m16-14 are ignored
	Program terminate	Provided $7168 \leq n \leq 8191$ current program level is terminated.	To be explained in Section 1.2.4 of this manual	19.0	The effect of this instruction on S will be explained in Section 1.2.4 of this manual.

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1. 2. 2.

This order code encompasses that of the MCD 920A computer. The following features are extensions to the 920A instruction code:

Block input

Block output

Preservation of Q by instructions 6 and 8.